

High Performance Molecular Simulation, Visualization, and Analysis on GPUs

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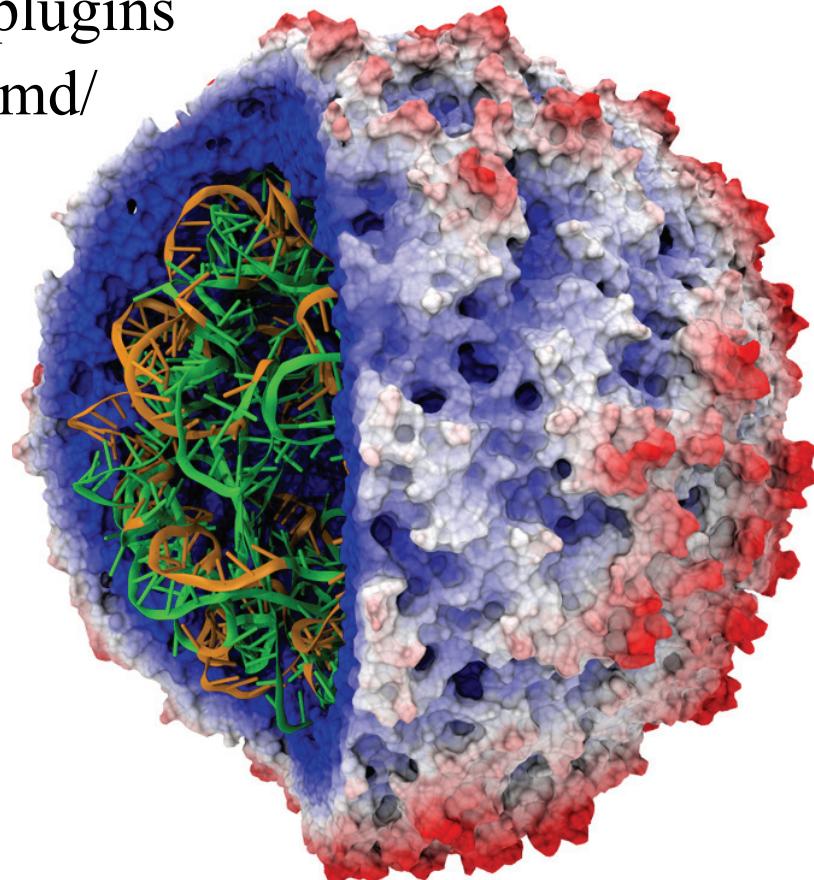
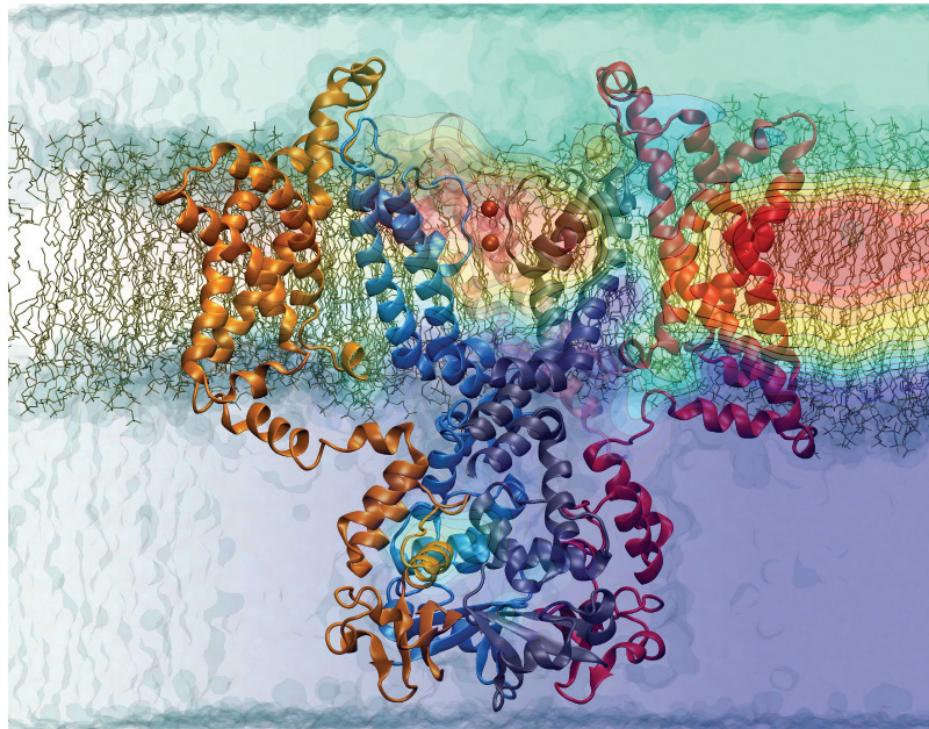
GPU Technology Conference

San Jose Convention Center,

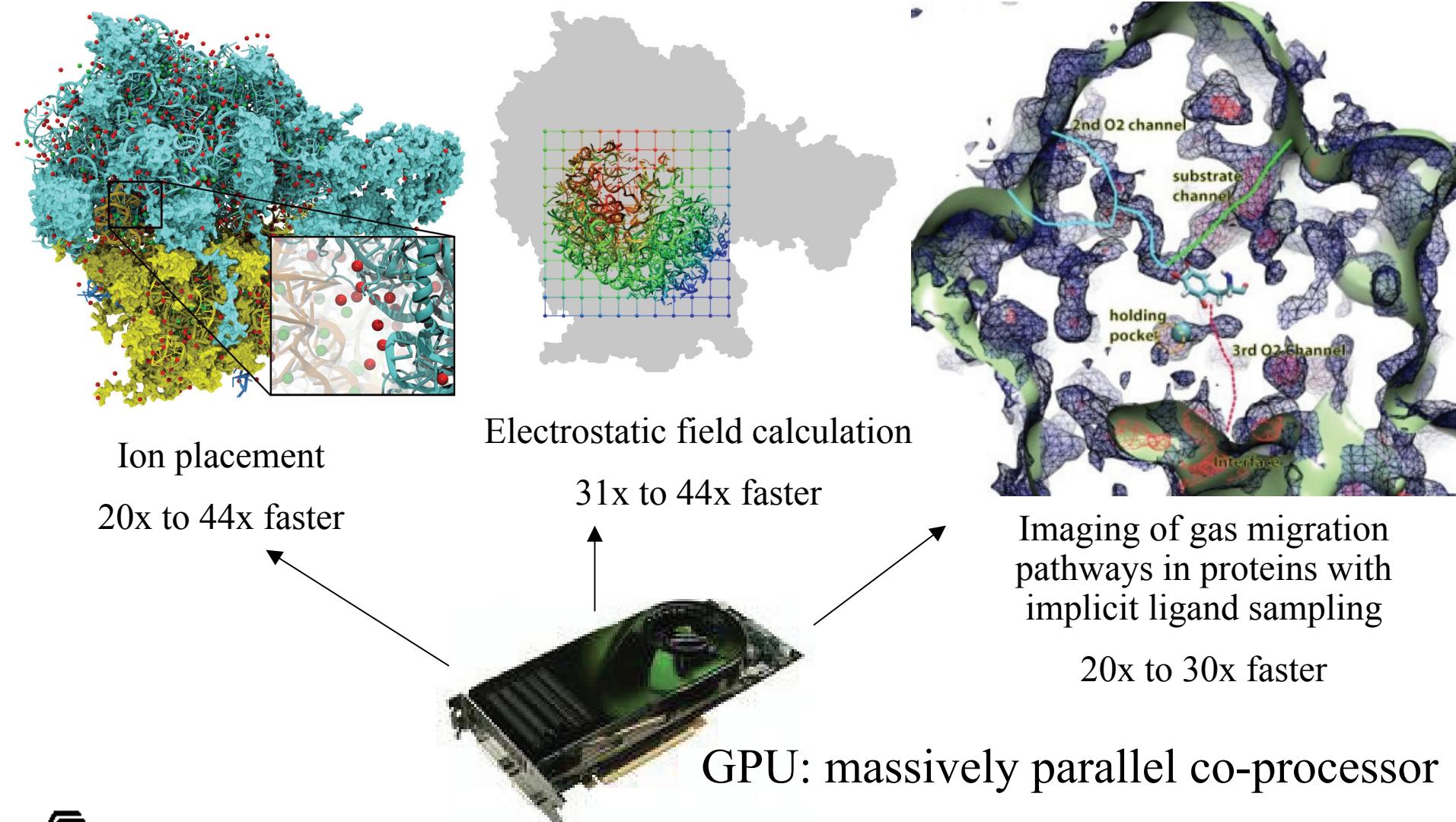
San Jose, CA, September 22, 2010

VMD – “Visual Molecular Dynamics”

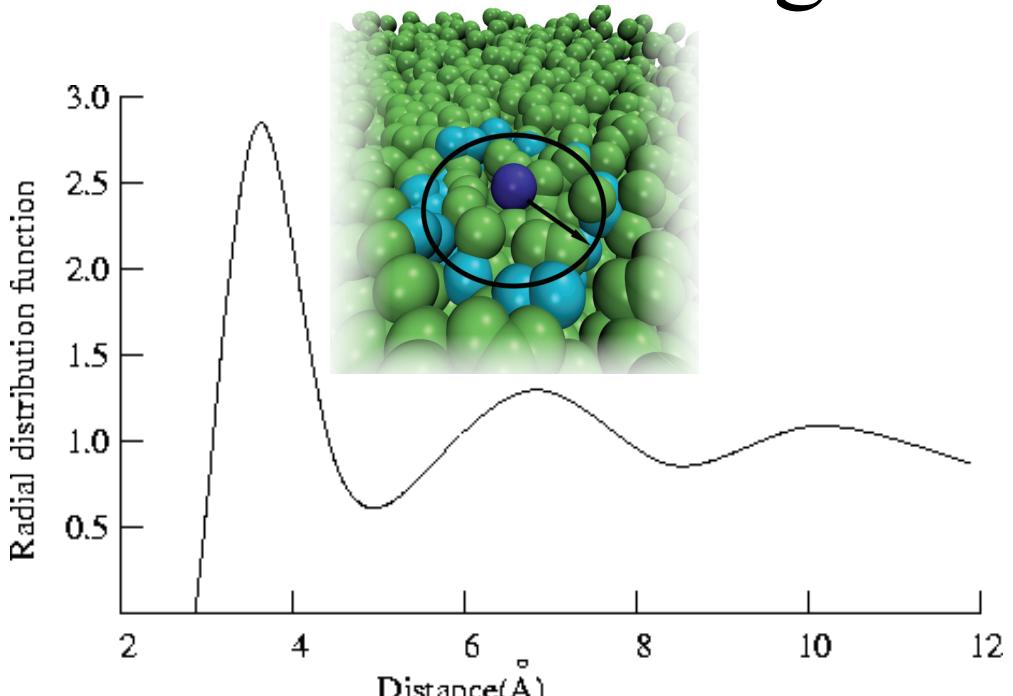
- Visualization and analysis of molecular dynamics simulations, sequence data, volumetric data, quantum chemistry simulations, particle systems, ...
- User extensible with scripting and plugins
- <http://www.ks.uiuc.edu/Research/vmd/>



CUDA Algorithms in VMD

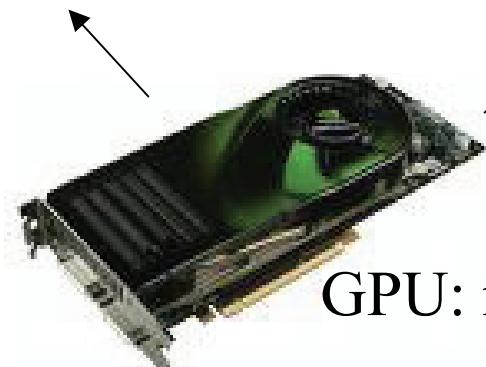


CUDA Algorithms in VMD

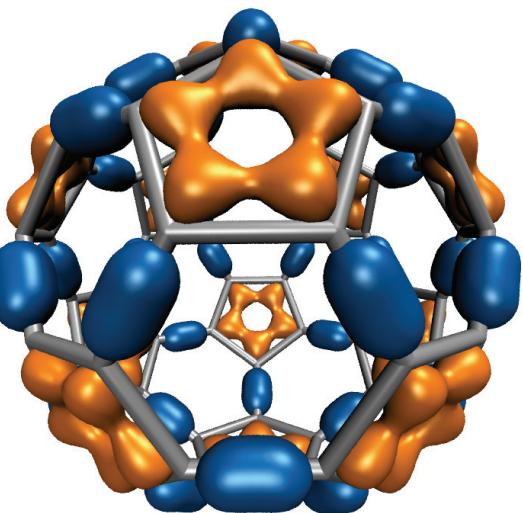


Radial distribution functions

30x to 92x faster



GPU: massively parallel co-processor



Molecular orbital
calculation and display

100x to 120x faster

Ongoing VMD GPU Development

- Development of new CUDA kernels for common molecular dynamics trajectory analysis tasks, faster surface renderings, and more...
- Support for CUDA in MPI-enabled builds of VMD for analysis runs on GPU clusters
- Updating existing CUDA kernels to take advantage of new hardware features on the latest NVIDIA “Fermi” GPUs
- Adaptation of some of our CUDA kernels to OpenCL to enable evaluation of JIT code generation techniques

Quantifying GPU Performance and Energy Efficiency in HPC Clusters

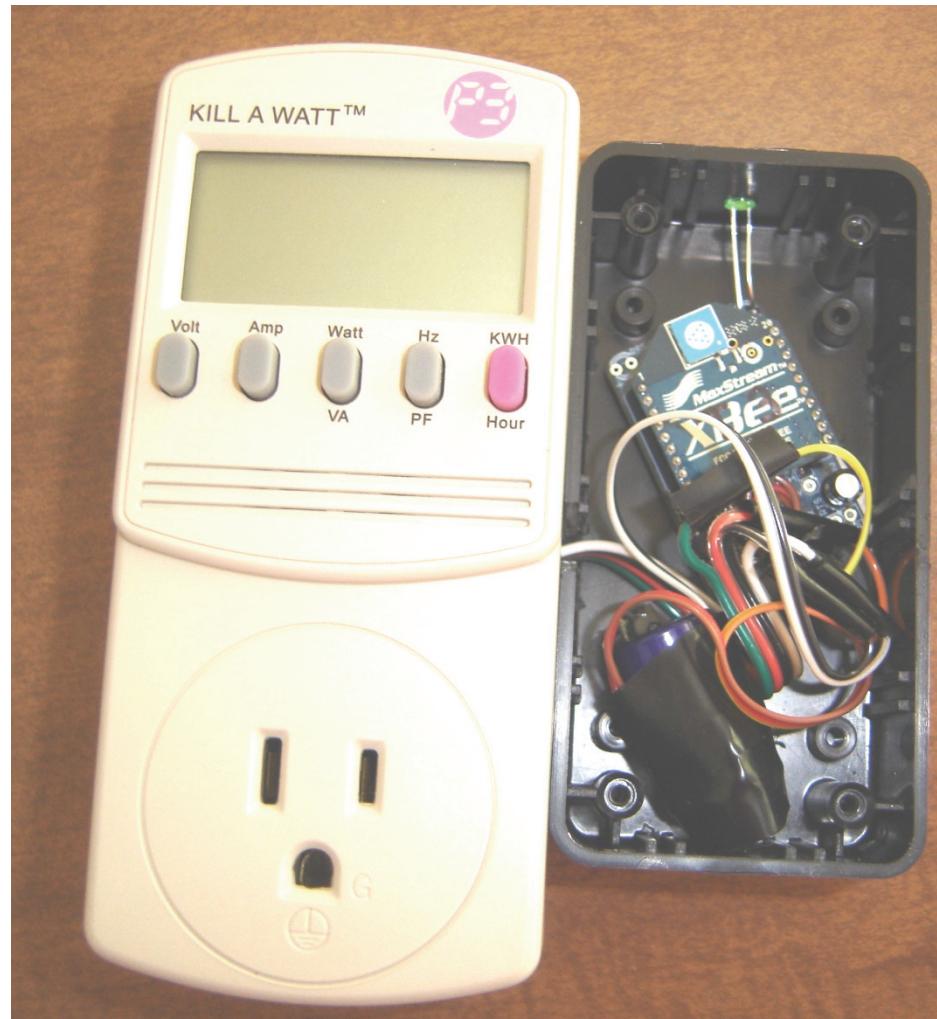
- NCSA “AC” Cluster
- Power monitoring hardware on one node and its attached Tesla S1070 (4 GPUs)
- Power monitoring logs recorded separately for host node and attached GPUs
- Logs associated with batch job IDs



- 32 HP XW9400 nodes
- 128 cores, 128 Tesla C1060 GPUs
- QDR Infiniband

Tweet-a-Watt

- Kill-a-watt power meter
- Xbee wireless transmitter
- Power, voltage, shunt sensing tapped from op amp
- Lower transmit rate to smooth power through large capacitor
- Readout software upload samples to local database
- We built 3 transmitter units and one Xbee receiver
- Currently integrated into AC cluster as power monitor



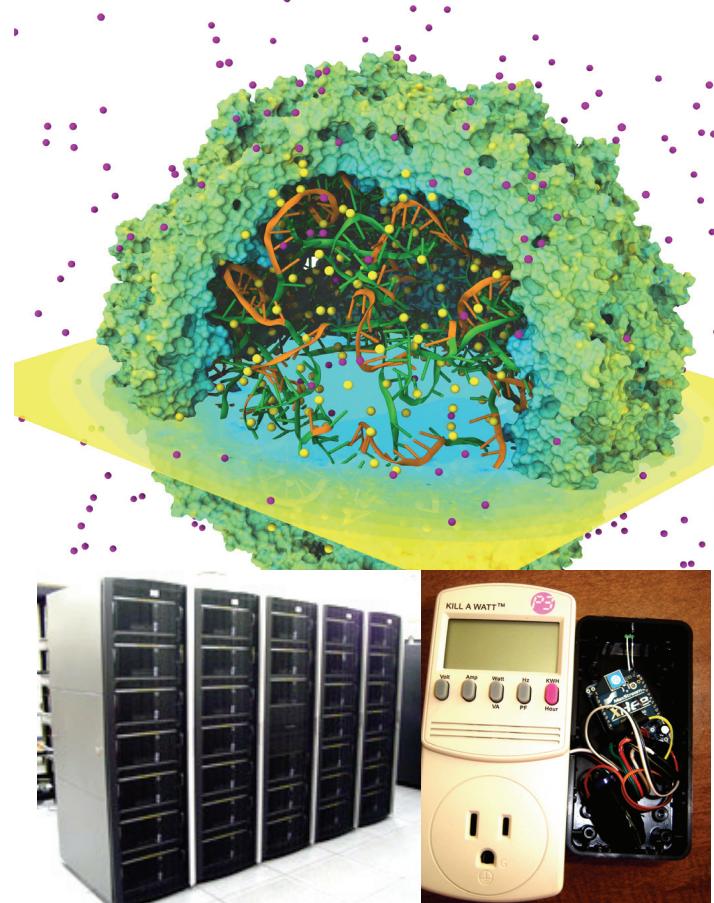
AC GPU Cluster Power Measurements

State	Host Peak (Watt)	Tesla Peak (Watt)	Host power factor (pf)	Tesla power factor (pf)
power off	4	10	.19	.31
pre-GPU use idle	173	178	.98	.96
after NVIDIA driver module unload/reload	173	178	.98	.96
after deviceQuery (idle)	173	365	.99	.99
GPU memtest #10 (stress)	269	745	.99	.99
VMD Multiply-add	268	598	.99	.99
NAMD GPU STMV	321	521	.97-1.0	.85-1.0

GPU Clusters for High Performance Computing. V. Kindratenko, J. Enos, G. Shi, M. Showerman, G. Arnold, J. Stone, J. Phillips, W. Hwu. *Workshop on Parallel Programming on Accelerator Clusters (PPAC)*, In Proceedings IEEE Cluster 2009, pp. 1-8, Aug. 2009.

Energy Efficient GPU Computing of Time-Averaged Electrostatics

- **1.5 hour job reduced to 3 min**
- Electrostatics of thousands of trajectory frames averaged
- Per-node power consumption on NCSA GPU cluster:
 - CPUs-only: 299 watts
 - CPUs+GPUs: 742 watts
- GPU Speedup: **25.5x**
- Power efficiency gain: **10.5x**



NCSA “AC” GPU cluster and Tweet-a-watt wireless power monitoring device

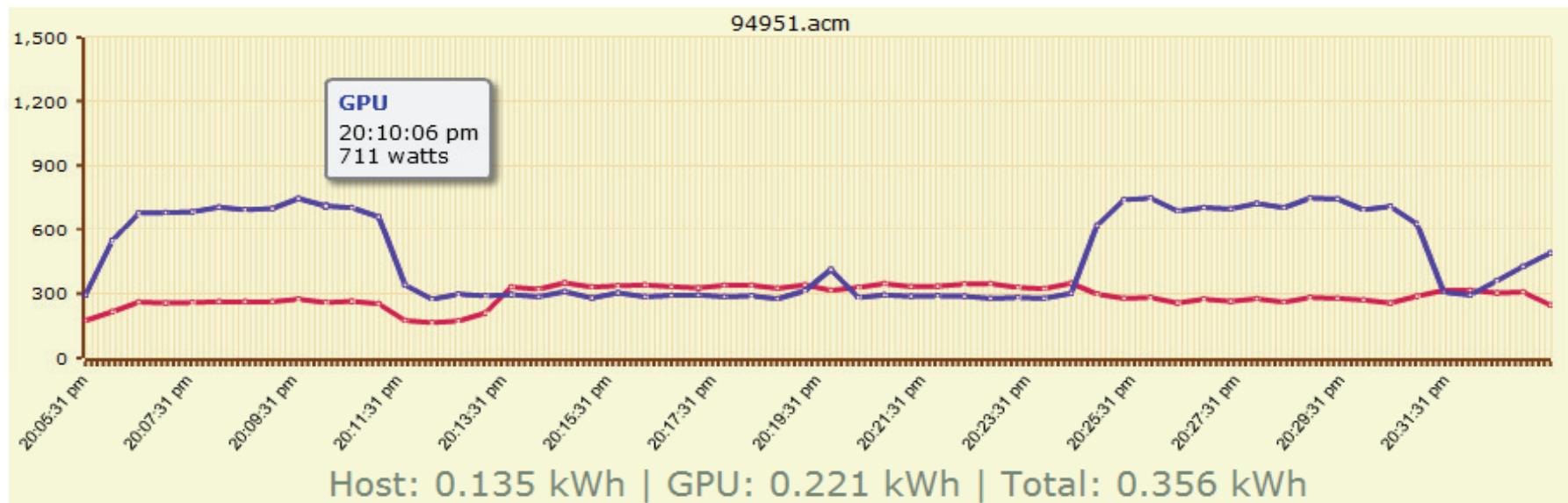
AC Cluster GPU Performance and Power Efficiency Results

Application	GPU speedup	Host watts	Host+GPU watts	Perf/watt gain
NAMD	6	316	681	2.8
VMD	25	299	742	10.5
MILC	20	225	555	8.1
QMCPACK	61	314	853	22.6

Quantifying the Impact of GPUs on Performance and Energy Efficiency in HPC Clusters. J. Enos, C. Steffen, J. Fullop, M. Showerman, G. Shi, K. Esler, V. Kindratenko, J. Stone, J. Phillips. *The Work in Progress in Green Computing*, 2010. In press.

Power Profiling: Example Log

AC Power Utilization



JSON Data

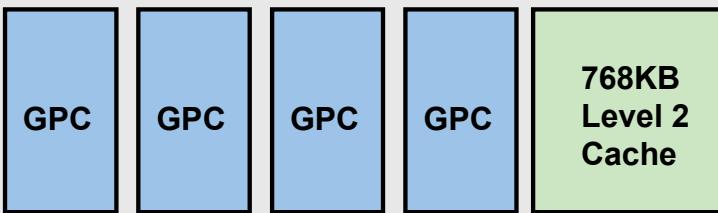
- Mouse-over value displays
- Under curve totals displayed
- If there is user interest, we may support calls to add custom tags from application

Fermi GPUs Bring Higher Performance and Easier Programming

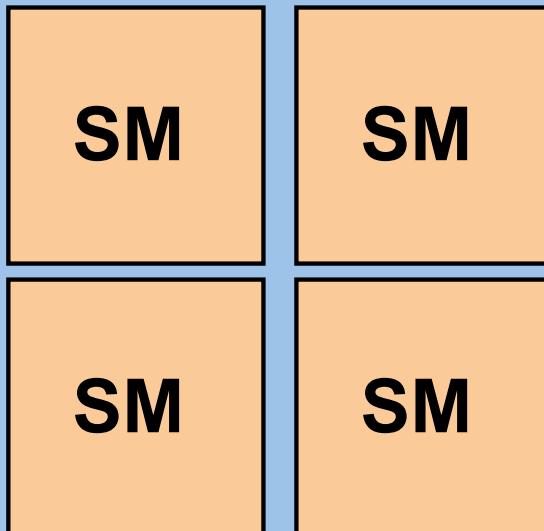
- NVIDIA's latest “Fermi” GPUs bring:
 - Greatly increased peak single- and double-precision arithmetic rates
 - Moderately increased global memory bandwidth
 - Increased capacity on-chip memory partitioned into shared memory and an L1 cache for global memory
 - Concurrent kernel execution
 - Bidirectional asynchronous host-device I/O
 - ECC memory, faster atomic ops, many others...

NVIDIA Fermi GPU

~3-6 GB DRAM Memory w/ ECC



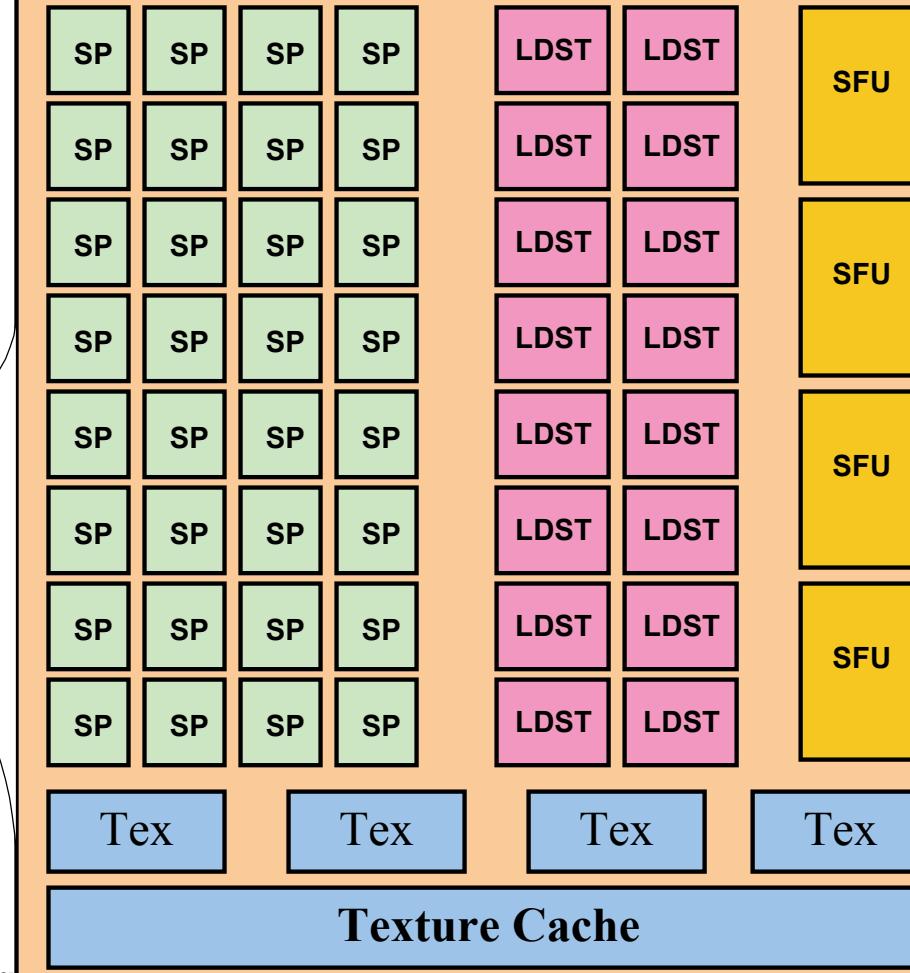
Graphics Processor Cluster



Streaming Multiprocessor

64KB Constant Cache

64 KB L1 Cache / Shared Memory

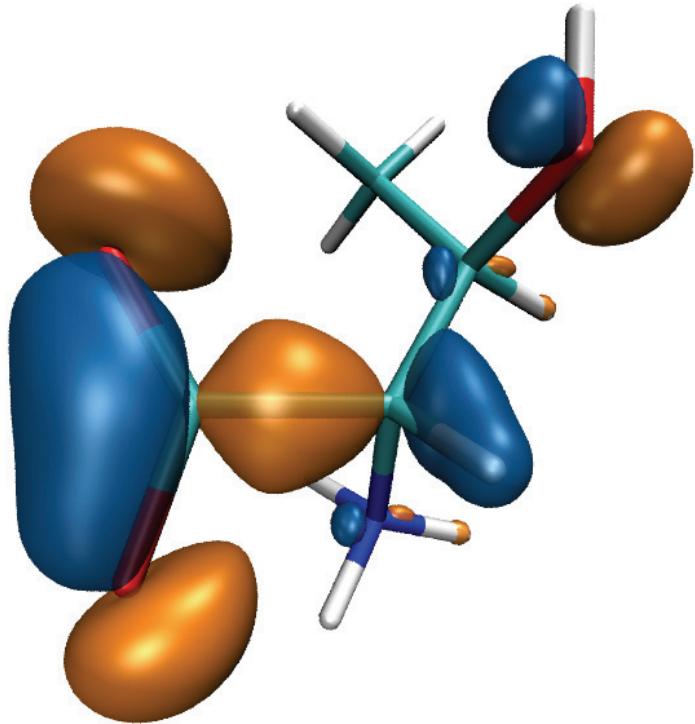


Early Experiences with Fermi

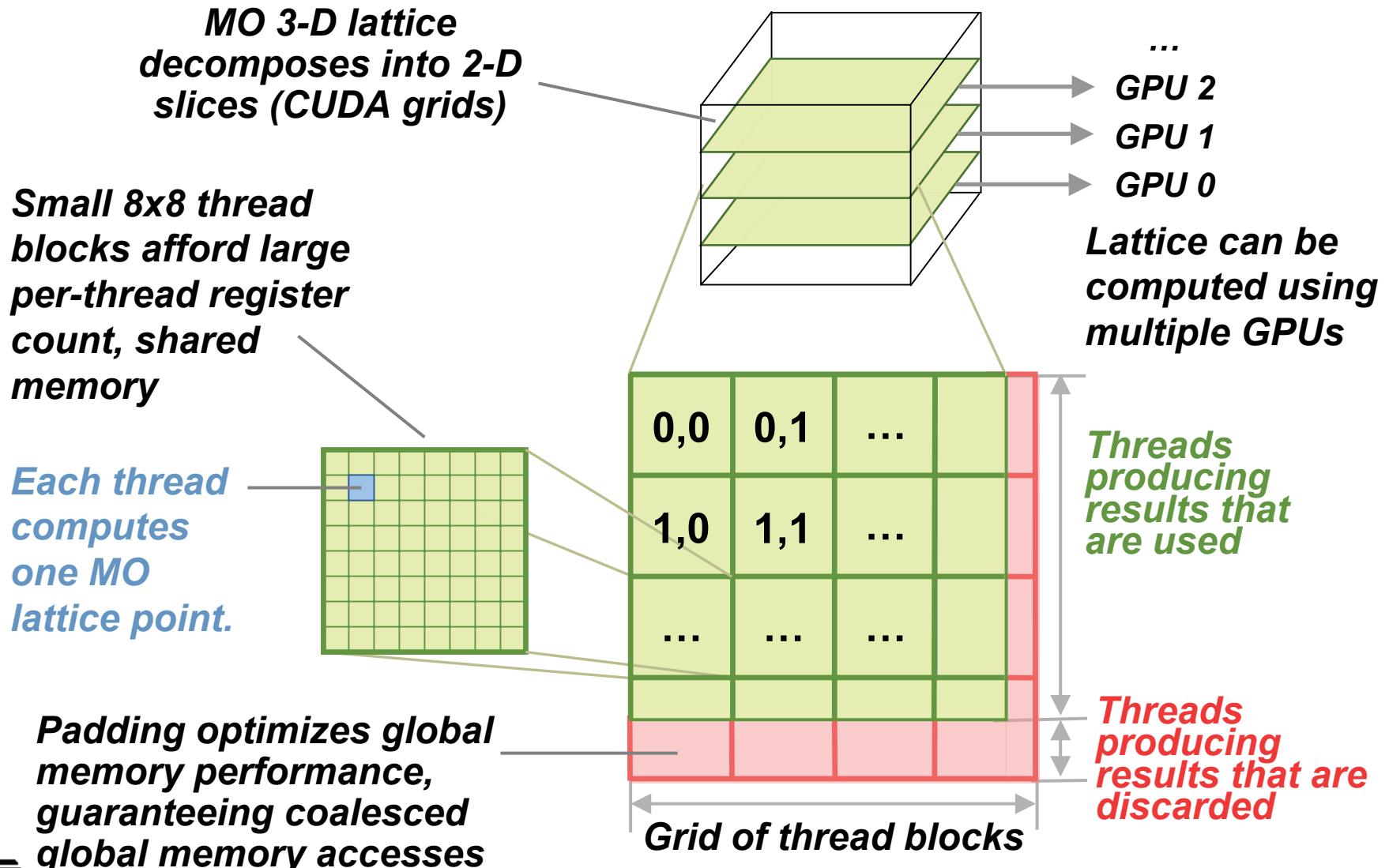
- The 2x single-precision and up to 8x double-precision arithmetic performance increases vs. GT200 cause more kernels to be memory bandwidth bound...
- ...unless they make effective use of the larger on-chip shared mem and L1 global memory cache to improve performance
- **Arithmetic is cheap, memory references are costly** (trend is certain to continue & intensify...)
- Register consumption and GPU “occupancy” are a bigger concern with Fermi than with GT200

Computing Molecular Orbitals

- Visualization of MOs aids in understanding the chemistry of molecular system
- Calculation of high resolution MO grids for display can require tens to hundreds of seconds on multi-core CPUs, even with the use of hand-coded SSE



MO GPU Parallel Decomposition



VMD MO GPU Kernel Snippet: Loading Tiles Into Shared Memory On-Demand

[... outer loop over atoms ...]

```
if ((prim_counter + (maxprim<<1)) >= SHAREDSIZE) {  
    prim_counter += sblock_prim_counter;  
    sblock_prim_counter = prim_counter & MEMCOAMASK;  
    s_basis_array[sidx] = basis_array[sblock_prim_counter + sidx];  
    s_basis_array[sidx + 64] = basis_array[sblock_prim_counter + sidx + 64];  
    s_basis_array[sidx + 128] = basis_array[sblock_prim_counter + sidx + 128];  
    s_basis_array[sidx + 192] = basis_array[sblock_prim_counter + sidx + 192];  
    prim_counter -= sblock_prim_counter;  
    __syncthreads();  
}  
  
for (prim=0; prim < maxprim; prim++) {  
    float exponent = s_basis_array[prim_counter];  
    float contract_coeff = s_basis_array[prim_counter + 1];  
    contracted_gto += contract_coeff * __expf(-exponent*dist2);  
    prim_counter += 2;  
}  
[... continue on to angular momenta loop ...]
```

Shared memory tiles:

- Tiles are checked and loaded, if necessary, immediately prior to entering key arithmetic loops

- Adds additional control overhead to loops, even with optimized implementation

VMD MO GPU Kernel Snippet: Fermi kernel based on L1 cache

```
[... outer loop over atoms ...]  
// loop over the shells/basis funcs belonging to this atom  
for (shell=0; shell < maxshell; shell++) {  
    float contracted_gto = 0.0f;  
    int maxprim  = shellinfo[(shell_counter<<4)      ];  
    int shell_type = shellinfo[(shell_counter<<4) + 1];  
    for (prim=0; prim < maxprim; prim++) {  
        float exponent      = basis_array[prim_counter      ];  
        float contract_coeff = basis_array[prim_counter + 1];  
        contracted_gto += contract_coeff * __expf(-  
            exponent*dist2);  
        prim_counter += 2;  
    }  
    [... continue on to angular momenta loop ...]
```

L1 cache:

- Simplifies code!
- Reduces control overhead
- Gracefully handles arbitrary-sized problems
- Matches performance of constant memory

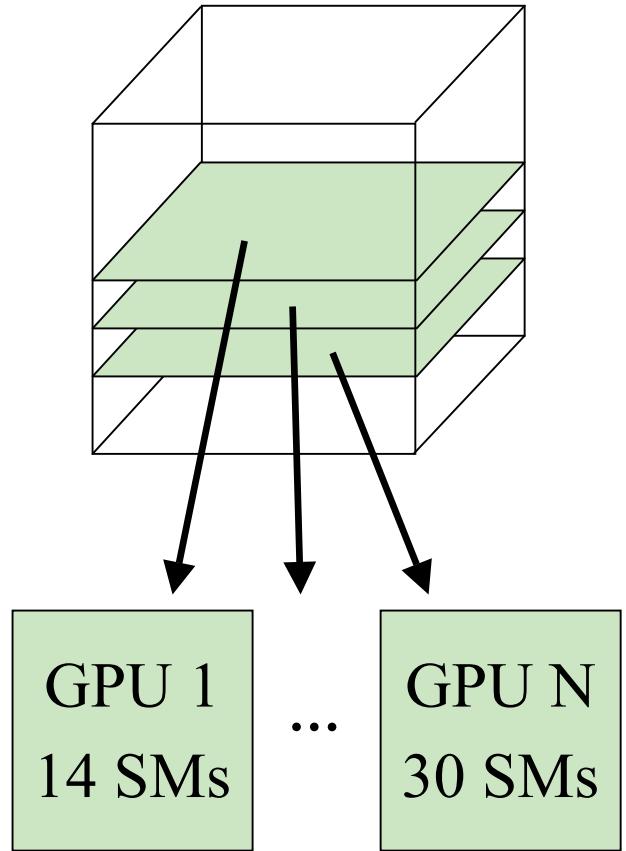
VMD Single-GPU Molecular Orbital Performance Results for C₆₀ on Fermi Intel X5550 CPU, GeForce GTX 480 GPU

Kernel	Cores/GPUs	Runtime (s)	Speedup
Xeon 5550 ICC-SSE	1	30.64	1.0
Xeon 5550 ICC-SSE	8	4.13	7.4
CUDA shared mem	1	0.37	83
CUDA L1-cache (16KB)	1	0.27	113
CUDA const-cache	1	0.26	117
CUDA const-cache, zero-copy	1	0.25	122

Fermi GPUs have caches: match perf. of hand-coded shared memory kernels. Zero-copy memory transfers improve overlap of computation and host-GPU I/Os.

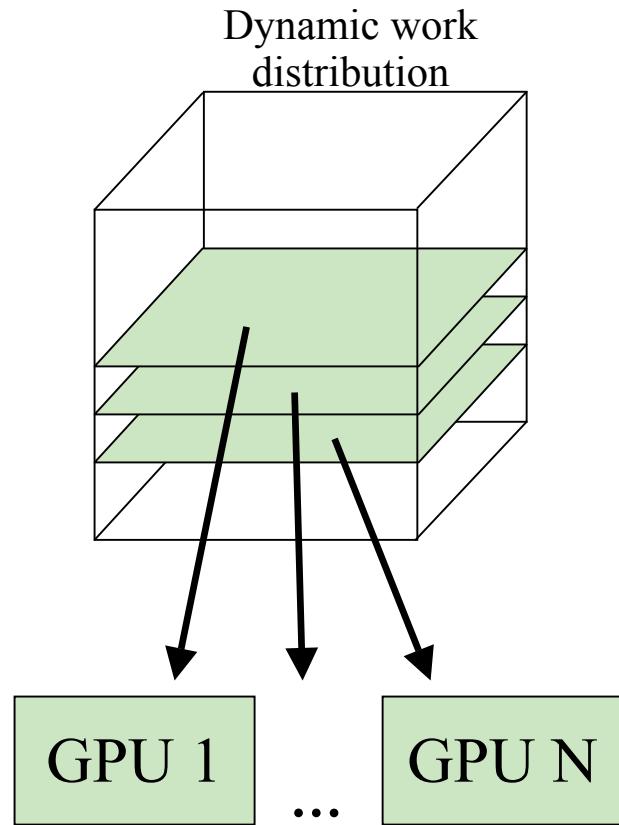
Multi-GPU Load Balance

- Many early CUDA codes assumed all GPUs were identical
- Host machines may contain a diversity of GPUs of varying capability (discrete, IGP, etc)
- Different GPU on-chip and global memory capacities may need different problem “tile” sizes
- Static decomposition works poorly for non-uniform workload, or diverse GPUs



Multi-GPU Dynamic Work Distribution

```
// Each GPU worker thread loops over  
// subset of work items...  
  
while (!threadpool_next_tile(&parms,  
    tilesize, &tile){  
  
    // Process one work item...  
  
    // Launch one CUDA kernel for each  
    // loop iteration taken...  
  
    // Shared iterator automatically  
    // balances load on GPUs  
  
}
```



Example Multi-GPU Latencies

4 C2050 GPUs, Intel Xeon 5550

6.3us	CUDA empty kernel (immediate return)
9.0us	Sleeping barrier primitive (non-spinning barrier that uses POSIX condition variables to prevent idle CPU consumption while workers wait at the barrier)
14.8us	pool wake, host fctn exec, sleep cycle (no CUDA)
30.6us	pool wake, 1x(tile fetch, simple CUDA kernel launch), sleep
1817.0us	pool wake, 100x(tile fetch, simple CUDA kernel launch), sleep

VMD Multi-GPU Molecular Orbital Performance Results for C₆₀

Intel X5550 CPU, 4x GeForce GTX 480 GPUs,

Kernel	Cores/GPUs	Runtime (s)	Speedup
Intel X5550-SSE	1	30.64	1.0
Intel X5550-SSE	8	4.13	7.4
GeForce GTX 480	1	0.255	120
GeForce GTX 480	2	0.136	225
GeForce GTX 480	3	0.098	312
GeForce GTX 480	4	0.081	378

Uses persistent thread pool to avoid GPU init overhead,
dynamic scheduler distributes work to GPUs

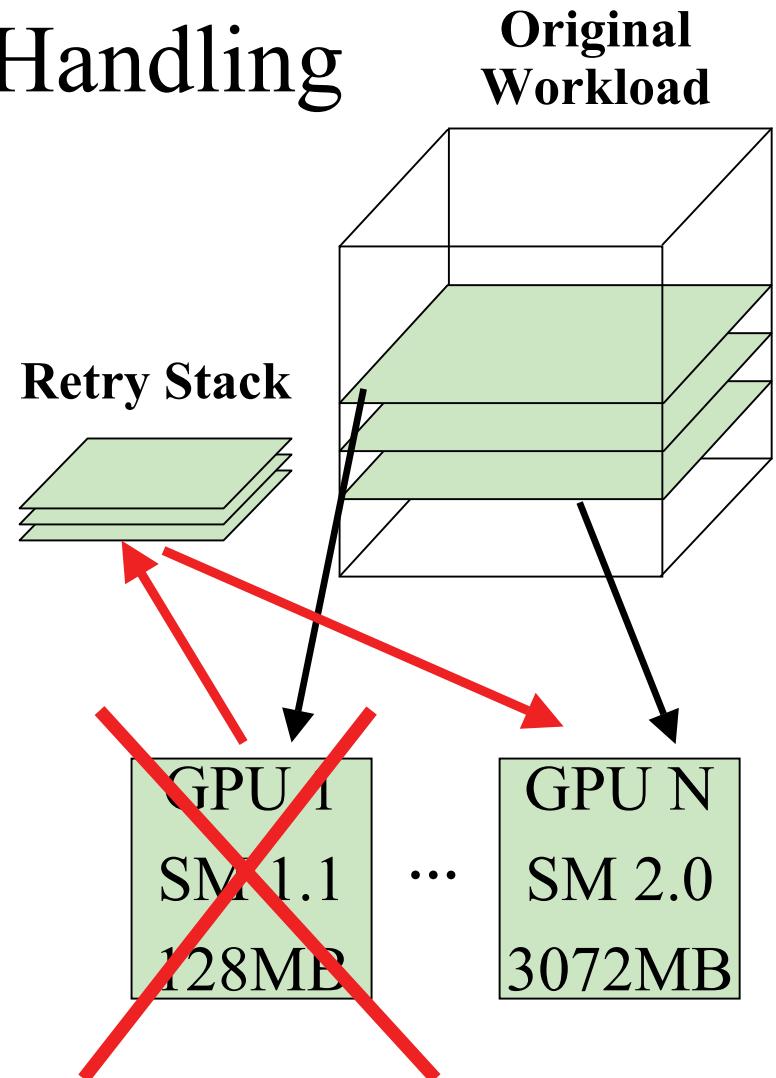
Molecular Orbital Dynamic Scheduling Performance with Heterogeneous GPUs

Kernel	Cores/GPUs	Runtime (s)	Speedup
Intel X5550-SSE	1	30.64	1.0
Quadro 5800	1	0.384	79
Tesla C2050	1	0.325	94
GeForce GTX 480	1	0.255	120
GeForce GTX 480 + Tesla C2050 + Quadro 5800	3	0.114	268 (91% of ideal perf)

Dynamic load balancing enables mixture of GPU generations, SM counts, and clock rates to perform well.

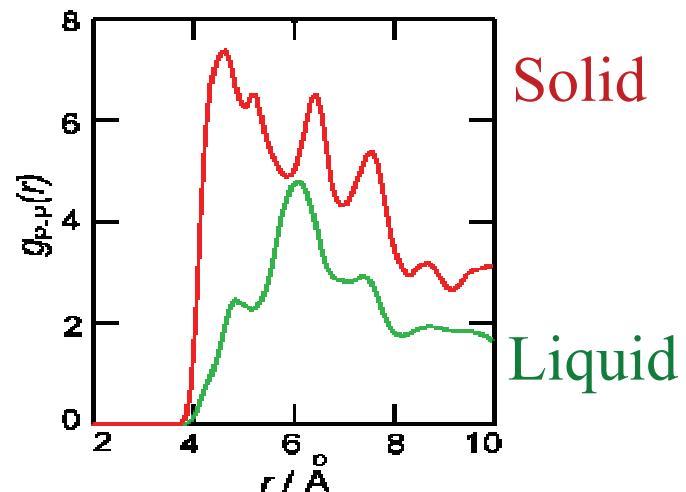
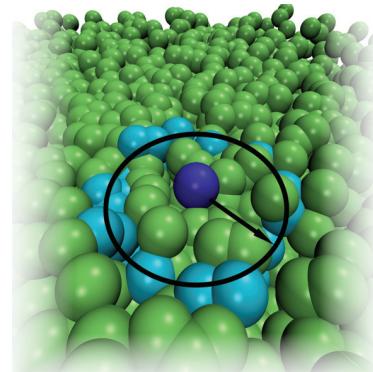
Multi-GPU Runtime Error/Exception Handling

- Competition for resources from other applications can cause runtime failures, e.g. GPU out of memory half way through an algorithm
- Handle exceptions, e.g. convergence failure, NaN result, insufficient compute capability/features
- Handle and/or reschedule failed tiles of work



Radial Distribution Function

- RDFs describes how atom density varies with distance
- Can be compared with experiments
- Shape indicates phase of matter: sharp peaks appear for solids, smoother for liquids
- Quadratic time complexity $O(N^2)$



Computing RDFs

- Compute distances for all pairs of atoms between two groups of atoms A and B
- A and B may be the same, or different
- Use nearest image convention for periodic systems
- Each pair distance is inserted into a histogram
- Histogram is normalized one of several ways depending on use, but usually according to the volume of the spherical shells associated with each histogram bin

Computing RDFs on CPUs

- Atom coordinates can be traversed in a strictly consecutive access pattern, yielding good cache utilization
- Since RDF histograms are usually small to moderate in size, they normally fit entirely in L2 cache
- CPUs compute the entire histogram in a **single pass**, regardless of the problem size or number of histogram bins

Parallel Histogramming on Multi-core CPUs

- Parallel updates to a single histogram bin creates a **potential output conflict**
- CPUs have atomic increment instructions, but they often take hundreds of clock cycles; unsuitable...
- For small numbers of CPU cores, it is best to **replicate** and **privatize** the histogram for each CPU thread, compute them independently, and combine the separate histograms in a final reduction step
- SSE can't be used effectively: lacks ability to “scatter” to memory (e.g. no *scatter-add*)

Computing RDFs on the GPU

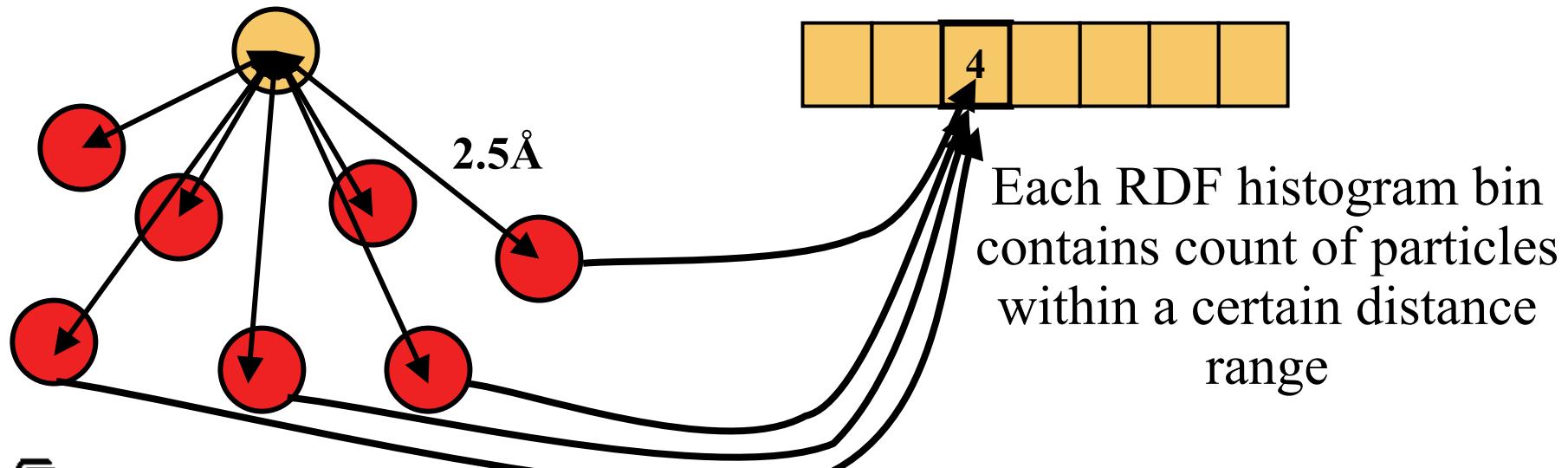
- Need tens of thousands of independent threads
- Each GPU thread computes one or more atom pair distances
- Performance is limited by the speed of histogramming
- Histograms are best stored in fast on-chip shared memory
- Small size of shared memory severely constrains the range of viable histogram update techniques
- **Fast CUDA implementation on Fermi: 30-92x faster than CPU**

Computing Atom Pair Distances on the GPU

- Memory access pattern is simple
- Primary consideration is **amplification of effective memory bandwidth**, through use of GPU on-chip shared memory, caches, and broadcast of data to multiple or all threads in a thread block

Radial Distribution Functions on GPUs

- Load blocks of atoms into shared memory and constant memory, compute periodic boundary conditions and atom-pair distances, all in parallel...
- Each thread computes all pair distances between its atom and all atoms in constant memory, incrementing the appropriate bin counter in the RDF histogram..



GPU Histogramming

- Tens of thousands of threads concurrently computing atom distance pairs...
- Far too many threads for a simple per-thread histogram privatization approach like CPU...
- Viable approach: **per-warp histograms**
- Fixed size shared memory limits histogram size that can be computed in a single pass
- Large histograms require **multiple passes**, but we can skip block pairs that are known not to contribute to a histogram window

Per-warp Histogram Approach

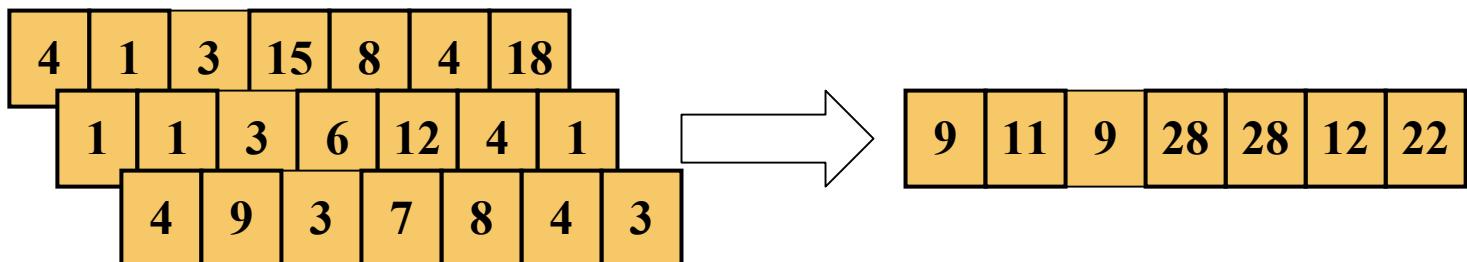
- Each warp maintains its own **private** histogram in on-chip shared memory
- Each thread in the warp computes an atom pair distance and updates a histogram bin in parallel
- Conflicting histogram bin updates are resolved using one of two schemes:
 - Shared memory write combining with thread-tagging technique (older hardware)
 - **atomicAdd()** to shared memory (new hardware)

RDF Inner Loops (abbreviated, xdist-only)

```
// loop over all atoms in constant memory
for (iblock=0; iblock<loopmax2; iblock+=3*NCUDABLOCKS*NBLOCK) {
    __syncthreads();
    for (i=0; i<3; i++) xyz[i[threadIdx.x + i*NBLOCK]]=pxi[iblock + i*NBLOCK]; // load coords...
    __syncthreads();
    for (joffset=0; joffset<loopmax; joffset+=3) {
        rxij=fabsf(xyz[i][idx3] - xyz[j][joffset]); // compute distance, PBC min image convention
        rxij2=celld.x - rxij;
        rxij=fminf(rxij, rxij2);
        rij=rxij*rxij;
        [...other distance components...]
        rij=sqrtf(rij + rxij*rxij);
        ibin=__float2int_rd((rij-rmin)*delr_inv);
        if (ibin<nbins && ibin>=0 && rij>rmin2) {
            atomicAdd(lhist+ibin, 1U);
        }
    } //joffset
} //iblock
```

Writing/Updating Histogram in Global Memory

- When thread block completes, add independent per-warp histograms together, and write to per-thread-block histogram in global memory
- Final reduction of all per-thread-block histograms stored in global memory

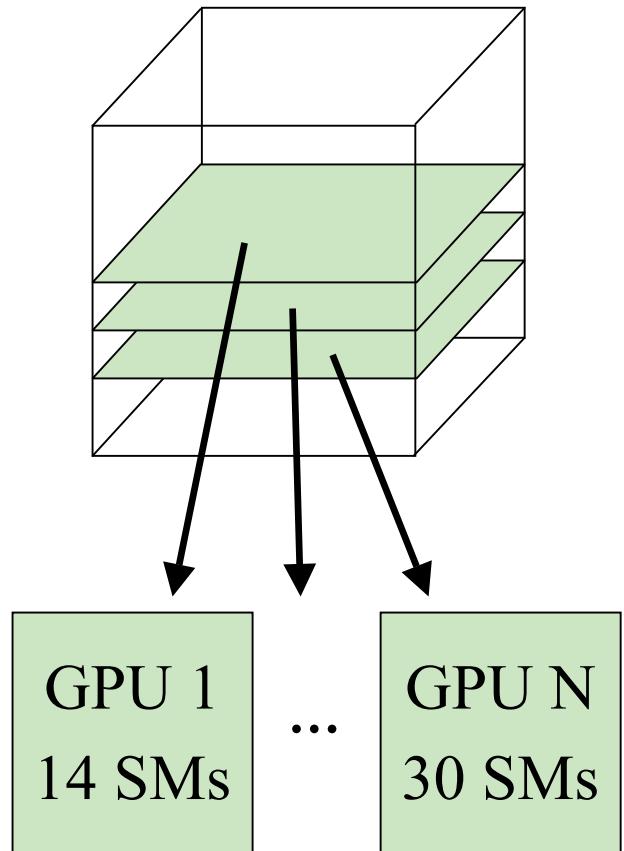


Preventing Integer Overflows

- Since all-pairs RDF calculation computes many billions of pair distances, we have to prevent integer overflow for the 32-bit histogram bin counters (supported by the atomicAdd() routine)
- We compute full RDF calculation in multiple kernel launches, so each kernel launch computes partial histogram
- Host routines read GPUs and increments large (e.g. long, or double) histogram counters in host memory after each kernel completes

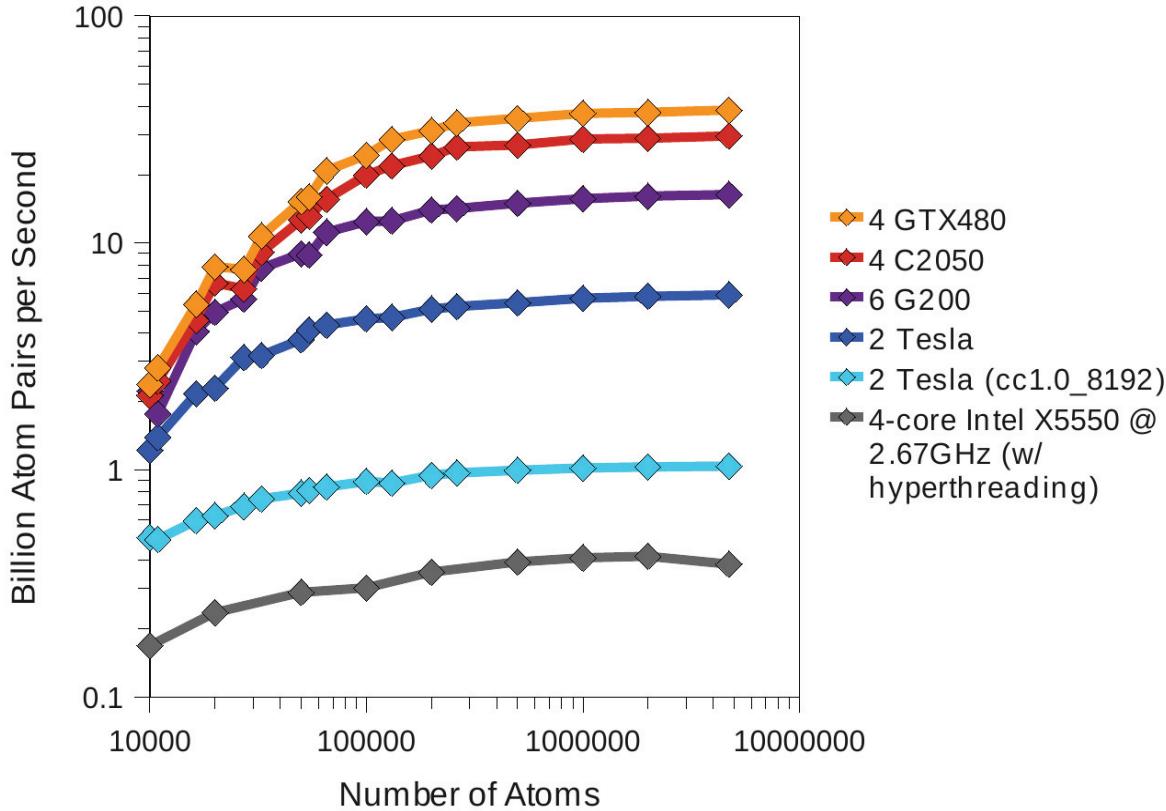
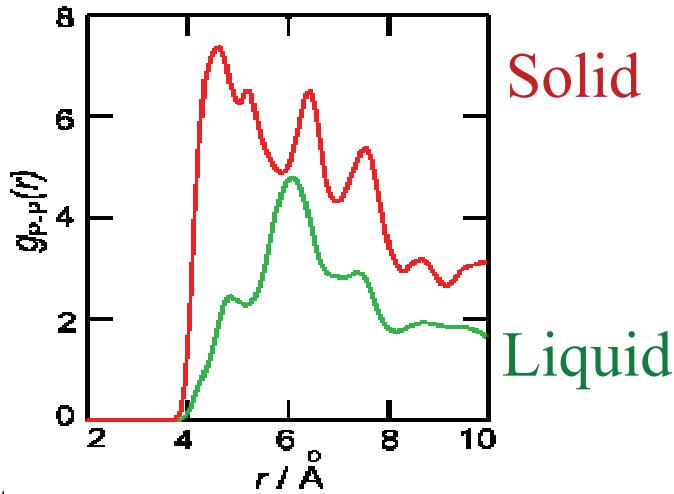
Multi-GPU RDF Calculation

- Distribute combinations of tiles of atoms and histogram regions to different GPUs
- Decomposed over two dimensions to obtain enough work units to balance GPU loads
- Each GPU computes its own histogram, and all results are combined for final histogram



Multi-GPU RDF Performance

- 4 NVIDIA GTX480 GPUs 30 to 92x faster than 4-core Intel X5550 CPU
- Fermi GPUs ~3x faster than GT200 GPUs: larger on-chip shared memory



Fast Analysis of Molecular Dynamics Trajectories with Graphics Processing Units – Radial Distribution Functions. B. Levine, J. Stone, and A. Kohlmeyer. 2010. (submitted)

Acknowledgements

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- NCSA Innovative Systems Lab
- The CUDA team at NVIDIA
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GPU Computing Publications

<http://www.ks.uiuc.edu/Research/gpu/>

- **Fast Analysis of Molecular Dynamics Trajectories with Graphics Processing Units – Radial Distribution Functions.** B. Levine, J. Stone, and A. Kohlmeyer. 2010. (submitted)
- **Quantifying the Impact of GPUs on Performance and Energy Efficiency in HPC Clusters.** J. Enos, C. Steffen, J. Fullop, M. Showerman, G. Shi, K. Esler, V. Kindratenko, J. Stone, J Phillips. *The Work in Progress in Green Computing*, 2010. In press.
- **GPU-accelerated molecular modeling coming of age.** J. Stone, D. Hardy, I. Ufimtsev, K. Schulten. *J. Molecular Graphics and Modeling*, 29:116-125, 2010.
- **OpenCL: A Parallel Programming Standard for Heterogeneous Computing.** J. Stone, D. Gohara, G. Shi. *Computing in Science and Engineering*, 12(3):66-73, 2010.
- **An Asymmetric Distributed Shared Memory Model for Heterogeneous Computing Systems.** I. Gelado, J. Stone, J. Cabezas, S. Patel, N. Navarro, W. Hwu. *ASPLOS '10: Proceedings of the 15th International Conference on Architectural Support for Programming Languages and Operating Systems*, pp. 347-358, 2010.

GPU Computing Publications

<http://www.ks.uiuc.edu/Research/gpu/>

- **Probing Biomolecular Machines with Graphics Processors.** J. Phillips, J. Stone. *Communications of the ACM*, 52(10):34-41, 2009.
- **GPU Clusters for High Performance Computing.** V. Kindratenko, J. Enos, G. Shi, M. Showerman, G. Arnold, J. Stone, J. Phillips, W. Hwu. *Workshop on Parallel Programming on Accelerator Clusters (PPAC)*, In Proceedings IEEE Cluster 2009, pp. 1-8, Aug. 2009.
- **Long time-scale simulations of in vivo diffusion using GPU hardware.** E. Roberts, J. Stone, L. Sepulveda, W. Hwu, Z. Luthey-Schulten. In *IPDPS'09: Proceedings of the 2009 IEEE International Symposium on Parallel & Distributed Computing*, pp. 1-8, 2009.
- **High Performance Computation and Interactive Display of Molecular Orbitals on GPUs and Multi-core CPUs.** J. Stone, J. Saam, D. Hardy, K. Vandivort, W. Hwu, K. Schulten, *2nd Workshop on General-Purpose Computation on Graphics Processing Units (GPGPU-2), ACM International Conference Proceeding Series*, volume 383, pp. 9-18, 2009.
- **Multilevel summation of electrostatic potentials using graphics processing units.** D. Hardy, J. Stone, K. Schulten. *J. Parallel Computing*, 35:164-177, 2009.

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<http://www.ks.uiuc.edu/Research/gpu/>

- **Adapting a message-driven parallel application to GPU-accelerated clusters.** J. Phillips, J. Stone, K. Schulten. *Proceedings of the 2008 ACM/IEEE Conference on Supercomputing*, IEEE Press, 2008.
- **GPU acceleration of cutoff pair potentials for molecular modeling applications.** C. Rodrigues, D. Hardy, J. Stone, K. Schulten, and W. Hwu. *Proceedings of the 2008 Conference On Computing Frontiers*, pp. 273-282, 2008.
- **GPU computing.** J. Owens, M. Houston, D. Luebke, S. Green, J. Stone, J. Phillips. *Proceedings of the IEEE*, 96:879-899, 2008.
- **Accelerating molecular modeling applications with graphics processors.** J. Stone, J. Phillips, P. Freddolino, D. Hardy, L. Trabuco, K. Schulten. *J. Comp. Chem.*, 28:2618-2640, 2007.
- **Continuous fluorescence microphotolysis and correlation spectroscopy.** A. Arkhipov, J. Hüve, M. Kahms, R. Peters, K. Schulten. *Biophysical Journal*, 93:4006-4017, 2007.